



## **Final project**

#### New visualization research or data analysis project

- **Research**: Pose problem, Implement creative solution
- **Data analysis**: Analyze dataset in depth & make a visual explainer

#### Deliverables

- **Research**: Implementation of solution
- Data analysis/explainer: Article with multiple interactive visualizations
- 6-8 page paper

#### Schedule

- Project proposal: Wed 2/19
- Design review and feedback: 3/9 and 3/11
- Final presentation: 3/16 (7-9pm) Location: TBD
- Final code and writeup: 3/18 11:59pm

#### Grading

- Groups of up to 3 people, graded individually
- Clearly report responsibilities of each member





# **Graphs and Trees**

**Graphs** Model relations among data Nodes and edges

## Trees

Graphs with hierarchical structure Connected graph with N-1 edges Nodes as parents and children



# **Spatial Layout**

### Primary concern – layout of nodes and edges

### Often (but not always) goal is to depict structure

- Connectivity, path-following
- Network distance
- Clustering
- Ordering (e.g., hierarchy level)

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## **Applications**

Tournaments Organization Charts Genealogy Diagramming (e.g., Visio) Biological Interactions (Genes, Proteins) Computer Networks Social Networks Simulation and Modeling Integrated Circuit Design

# Topics

Tree Layout Network Layout Sugiyama-Style Layout Force-Directed Layout Alternatives to Network Layout Matrix Diagrams Attribute-Drive Layout







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## Separate breadth & depth in 2D Focus on single path at a time

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# **Basic Recursive Approach**

Repeatedly divide space for subtrees by leaf count

- Breadth of tree along one dimension
- Depth along the other dimension





## **Basic Recursive Approach**

Repeatedly divide space for subtrees by leaf count

- Breadth of tree along one dimension
- Depth along the other dimension

Problem: Exponential growth of breadth



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## **Reingold & Tilford's Tidier Layout**



**Goal:** maximize density and symmetry.

Originally for binary trees, extended by Walker to cover general case.

This extension was corrected by Buchheim et al. to achieve a linear time algorithm

## **Reingold-Tilford Layout**

## **Design concerns**

Clearly encode depth level No edge crossings Isomorphic subtrees drawn identically Ordering and symmetry preserved Compact layout (don 't waste space)

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# **Radial Layout**



Node-link diagram in polar coords

Radius encodes depth root at center

Angular sectors assigned to subtrees (recursive approach)

Reingold-Tilford approach can also be applied here

## **Problems with Node-Link Diagrams**

## Scale

Tree breadth often grows exponentially Even with tidier layout, quickly run out of space

## **Possible solutions**

Filtering Focus+Context Scrolling or Panning Zooming Aggregation





# Hyperbolic Layout





Layout in hyperbolic space, then project on to Euclidean plane

Why? Like tree breadth, the hyperbolic plane expands exponentially

Also computable in 3D, projected into a sphere





## **Enclosure Diagrams**

Encode structure using spatial enclosure Popularly known as TreeMaps

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## **Benefits**

Provides a single view of an entire tree Easier to spot large/small nodes

## Problems

Difficult to accurately read depth

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# **Circle Packing Layout**



# **Circle Packing Layout**



Nodes represented as sized circles

Nesting to show parent-child relationships

**Problems**: Inefficient use of space Parent size misleading

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## Treemaps

Hierarchy visualization that emphasizes values of nodes via area encoding

Partition 2D space such that leaf nodes have sizes proportional to data values

First layout algorithms proposed by <u>Shneiderman et al. in</u> <u>1990</u>, with focus on showing file sizes on a hard drive





# Squarified Treemaps [Bruls 00]

Greedy optimization for objective of square rectangles Slice/dice within siblings; alternate whenever ratio worsens



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# Why Squares

### Posited Benefits of 1:1 Aspect Ratios

- 1. Minimize perimeter, reducing border ink.
- 2. Easier to select with a mouse cursor. Validated by empirical research & Fitt's Law!
- 3. Similar aspect ratios are easier to compare. Seems intuitive, but is this true?













## Voronoi Treemaps [Balzer 05]

Treemaps with arbitrary polygonal shape and boundary

Uses iterative, eighted Voronoi tessellations to achieve cells with valueproportional areas













## **Spanning Tree Layout**

Many graphs are tree-like or have useful spanning trees

Websites, Social Networks

Use tree layout on spanning tree of graph Trees created by BFS / DFS Min/max spanning trees

Fast tree layouts allow graph layouts to be recalculated at interactive rates

Heuristics may further improve layout



# Sugiyama-style graph layout



# Sugiyama-style graph layout



Reverse some edges to remove cycles Assign nodes to hierarchy layers → Longest path layering Create dummy nodes to "fill in" missing layers Arrange nodes within layer, minimize edge crossings Route edges – layout splines if needed





























## **Force-Directed Layout**

Nodes = charged particles with air resistance Edges = springs  $F = q_i * q_j / d_{ij^2}$   $F = -b * v_i$  $F = k * (L - d_{ij})$ 

## D3's force layout uses velocity Verlet integration

Assume uniform mass *m* and timestep  $\Delta t$ :  $F = ma \rightarrow F = a \rightarrow F = \Delta v / \Delta t \rightarrow F = \Delta v$ Forces simplify to velocity offsets!

Repeatedly calculate forces, update node positions Naïve approach O(N<sup>2</sup>) Speed up to O(N log N) using quadtree or k-d tree Numerical integration of forces at each time step















