# Nełwork Layout 

## Maneesh Agrawala

## CS 448B: Visualization Fall 2021

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## Reading Response Questions/Thoughts

For the final project, do you have a recommendation of a place to go to view other data visualization research papers that conducted user studies?

As animations contain more and more data, is it possible that we can overload or overstimulate the user? Can animations be harmful by being too distracting? If so, how can we safeguard our designs to make sure they don't cause this overstimulation?

Is there a more formal or mathematical rule set governing which colors to use to highlight information, and which to contrast? Or is it mostly a combination of multiple factors that you need to see to know? In a similar vein, do colors need to be different in shade as well as color for black and white printing? How do we know to vary transparency with color or just color?

How seriously should we take self-reported stated preferences when evaluating the strength of a visualization? How much should we weight user's expressed preference relative to usability, learning, and recall data when evaluating the efficacy of a visualization?

## Last Time: Animation Understanding Motion

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## How does it work?



Two-cylinder Stirling engine
http://www.keveney.com/Vstirling.html

## Problems [Tversky 02]

## Difficulties in understanding animation

- Difficult to estimate paths and trajectories
- Motion is fleeting and transient
- Cannot simultaneously attend to multiple motions
- Trying to parse motion into events, actions and behaviors
- Misunderstanding and wrongly inferring causality
- Anthropomorphizing physical motion may cause confusion or lead to incorrect conclusions

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## Solution I: Break into static steps



## Two-cylinder Stirling engine <br> http://www.keveney.com/Vstirling.html

## Challenges

## Choosing the set of steps

$\square$ How to segment process into steps?

- Note: Steps often shown sequentially for clarity, rather than showing everything simultaneously


## Tversky suggests

- Coarse level - segment based on objects
- Finer level - segment based on actions
- Static depictions often do not show finer level segmentation

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## Animated Transitions in Statistical Graphics



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## Log Transform




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Sorting



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Filtering



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Change Data Dimensions


Change Data + Encodings


## Change Encodings + Axis Scales



## Data Graphics \& Transitions

| Category | Sales | Profit |
| :--- | ---: | ---: |
| A | 11 | 7 |
| B | 13 | 10 |
| C | 12 | 6 |
| D | 8 | 5 |
| E | 3 | 1 |



Change selected data
dimensions or encodings
Animation to
municate changes?

| Category | Sales | Profit |
| :--- | ---: | ---: |
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# Animated Transitions in Statistical Data Graphics 

Jeffrey Heer<br>George G. Robertson

## Miorosoft <br> Research

## Study Conclusions

Appropriate animation improves graphical perception
Use simple staged transitions, but doing one thing at a time not always best

Axis re-scaling hampers perception
Avoid if possible (use common scale)
Maintain landmarks better (delay fade out of gridlines)
Subjects preferred animated transitions

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## Implementing Animation

## Animation Approaches

## Frame-based Animation

Redraw scene at regular interval (e.g., 16 ms )
Developer defines the redraw function

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## Frame-based Animation


circle $(15,15)$

circle $(25,25)$


1
2


| 1 | 2 | 3 | 4 |
| :--- | :--- | :--- | :--- |

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## Animation Approaches

Frame-based Animation<br>Redraw scene at regular interval (e.g., 16 ms )<br>Developer defines the redraw function

## Animation Approaches

Frame-based Animation
Redraw scene at regular interval (e.g., 16 ms )
Developer defines the redraw function
Transition-based Animation (Hudson \& Stasko '93)
Specify property value, duration \& easing (łweening)
Typically computed via interpolation

$$
\text { step }(\text { fraction })\left\{\mathrm{x}_{\text {now }}=\mathrm{x}_{\text {start }}+\text { fraction * }\left(\mathrm{Xend}_{\text {end }}-\mathrm{x}_{\text {start }}\right) ;\right\}
$$

Timing \& redraw managed by UI toolkit

## Transition-based Animation

from: $(10,10)$ to: $(25,25)$ duration: 3sec

$$
d x=25-10
$$

$$
x=10+(t / 3) * d x
$$


$x=10+(t / 3) * d x$

$x=10+(t / 3) * d x$


## Transition-based Animation

from: $(10,10)$ to: $(25,25)$ duration: 3sec
Toolkit handles frame-by-frame updates


## D3 Transitions

Any d3 selection can be used to drive animation.

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// Select SVG rectangles and bind them to data values
var bars = svg.selectAll("rect.bars").data(values);

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## Any d3 selection can be used to drive animation.

// Select SVG rectangles and bind them to data values.
var bars = svg.selectAll("rect.bars").data(values);
// Static transition: update position and color of bars.
bars
.attr("x", (d) => xScale(d.foo))
.attr("y", (d) => yScale(d.bar))
.style("fill", (d) => colorScale(d.baz));

## D3 Transitions

Any d3 selection can be used to drive animation.
// Select SVG rectangles and bind them to data values.
var bars = svg.selectAll("rect.bars").data(values);
// Animated transition: interpolate to target values using default timing
bars.transition()

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// Animation is implicitly queued to run!

## D3 Transitions, Continued

## bars.transition()

```
.duration(500) // animation duration in ms
    .delay(0) // onset delay in ms
    .ease(d3.easeBounce) // set easing (or "pacing") style
    .attr("x",(d) => xScale(d.foo))
    ...
```


## D3 Transitions, Continued

bars.transition()
.duration(500) // animation duration in ms
.delay(0) // onset delay in ms
.ease(d3.easeBounce) // set easing (or "pacing") style .attr("x", (d) => xScale(d.foo))
...
bars.exit().transition() // animate elements leaving display
.style("opacity", 0) // fade out to fully transparent
.remove(); // remove from DOM upon completion

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## Easing Functions

Goals: stylize animation, improve perception.
Basic idea is to warp time: as duration goes from start (0\%) to end (100\%), dynamically adjust the interpolation fraction using an easing function.


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## Summary

Animation is a salient visual phenomenon
Attention, object constancy, causality, timing
For processes, step-by-step static images may be preferable
For transitions, animation has some benefits, but consider task and timing

## Announcements

## Final project

Data analysis/explainer or conduct research

- Data analysis: Analyze dataset in depth \& make a visual explainer
- Research: Pose problem, Implement creative solution

Deliverables

- Data analysis/explainer: Article with multiple different interactive visualizations
- Research: Implementation of solution and web-based demo if possible
- Short video (2 $\mathbf{~ m i n}$ ) demoing and explaining the project


## Schedule

- Project proposal: Wed 11/3
- Design Review and Feedback: $10^{\text {th }}$ week of quarter
- Final code and video: Fri 12/10 11:59pm


## Grading

- Groups of up to 3 people, graded individually
- Clearly report responsibilities of each member


## Nełwork Layouł



## Graphs and Trees

Graphs
Model relations among data
Nodes and edges


## Trees

Graphs with hierarchical structure
Connected graph with N -1 edges Nodes as parents and children


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## Tree Layouł



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## Indentation



## Visualizing Large Hierarchies



Indented Layout

## Single-Focus (Accordion) List



Separate breadth \& depth in 2D Focus on single path at a time

## Node-Link Diagrams

Nodes distributed in space, connected by lines
Use 2D space to break apart breadth and depth
Space used to communicate hierarchical orientation
Typically towards authority or generality


## Basic Recursive Approach

Repeatedly divide space for subtrees by leaf count

- Breadth of tree along one dimension
- Depth along the other dimension



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Problem: Exponential growth of breadith


## Reingold \& Tilford's Tidier Layout



Goal: maximize density and symmetry.

Originally for binary trees, extended by Walker to cover general case.

This extension was corrected by Buchheim et al. to achieve a linear time algorithm

## Reingold-Tilford Layout

Design concerns
Clearly encode depth level
No edge crossings
Isomorphic subtrees drawn identically
Ordering and symmetry preserved
Compact layout (don 't waste space)

## Reingold-Tilford Algorithm

Initial bottom-up (postorder) tree traversal

- Set $y$-coordinate based on depth
- Initialize $x$-coordinate to zero

At each parent node, merge left and right subtrees

- Shift right subtree as close as possible to left
- Computed efficiently by maintaining subtree contours
- Center parent nodes above children
- Record "Shiff" in position offset for right subtree

Find top-down (preorder) traversal to set x-coordinates

- Sum aggregated shift

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## Reingold-Tilford Algorithm



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## Radial Layouł



Node-link diagram in polar coords
Radius encodes depth root at center
Angular sectors assigned to subtrees (recursive approach)

Reingold-Tilford approach can also be applied here

## Problems with Node-Link Diagrams

## Scale

Tree breadth often grows exponentially
Even with fidier layout, quickly run out of space

## Possible solutions

Filtering
Focus+Context
Scrolling or Panning
Zooming
Aggregation

## Visualizing Large Hierarchies



Indented Layout


Reingold-Tilford Layout

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## Hyperbolic Layout



Layout in hyperbolic space, then project on to Euclidean plane

Why? Like tree breadth, the hyperbolic plane expands exponentially

Also computable in 3D, projected into a sphere

## Degree-of-Interest Trees [AVI 04]



Space-constrained, multi-focal tree layout

## Degree-of-Interest Trees



Cull "un-interesting" nodes on a per block basis until all blocks on a level fit within bounds
Center child blocks under parents

## Enclosure Diagrams

Encode structure using spatial enclosure Popularly known as TreeMaps

## Benefits



Provides a single view of an entire tree
Easier to spot large/small nodes
Problems
Difficult to accurately read depth

## Circle Packing Layout



Nodes represented as sized circles
Nesting to show parent-child relationships

Problems:

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## Circle Packing Layouł



Nodes represented as sized circles
Nesting to show parent-child relationships

Problems:
Inefficient use of space
Parent size misleading

## Treemaps

Hierarchy visualization that emphasizes values of nodes via area encoding

Partition 2D space such that leaf nodes have sizes proportional to data values

First layout algorithms proposed by Shneiderman et al. in 1990, with focus on showing file sizes on a hard drive


Slice \& Dice layout: Alternate horizontal / vertical partitions.


Squarifed layout: Try to produce square (1:1) aspect ratios

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## Squarified Treemaps [Bruls 00]

Greedy optimization for objective of square rectangles Slice/dice within siblings; alternate whenever ratio worsens


## Why Squares

Posited Benefits of 1:1 Aspect Ratios

1. Minimize perimeter, reducing border ink.
2. Easier to select with a mouse cursor.

Validated by empirical research \& Fitt's Law!
3. Similar aspect ratios are easier to compare.

Seems intuitive, but is this true?

## Error vs. Aspect Rctio [Kong 10]



1. Comparison of squares has higher error!
2. Squarify works because it fails to meet its objective?

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Seems intuitive, but is this true?
Extreme ratios \& squares-only more inaccurate.
Balanced ratios better? Target golden ratio?

